
ECE 4514

Digital Design II

Spring 2008

Lecture 21:

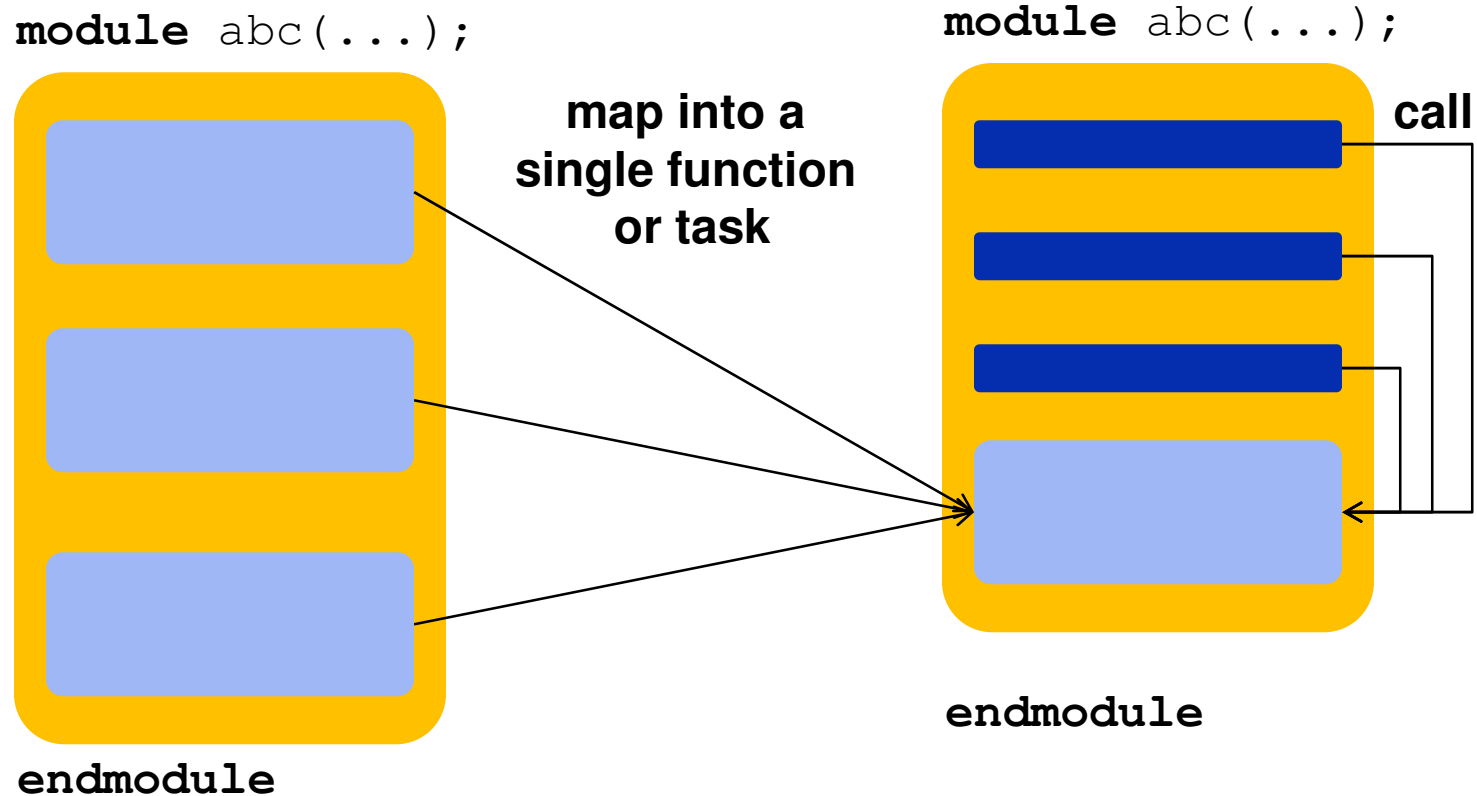
Functions and Tasks

A Language Lecture

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Tasks and Functions

- Help with the reusability of writing code



Tasks vs Functions

- Tasks, as well as functions, are *behavioral* constructs.
- They support expressions, if-then, case, loop statements
- They are local to a module, and called from within initial and always blocks.

Tasks

- ❑ Can contain event control (@, #, wait)
- ❑ May execute in non-zero simulation time
- ❑ Can have zero or more arguments of type in, out, inout

Functions

- ❑ Cannot contain event control (@, #, wait)
- ❑ Must execute in zero simulation time
- ❑ Must have at least in input argument, and has a single output argument

Functions

```
module andorgate(q, a, b);
    output [1:0] q;
    input      a, b;

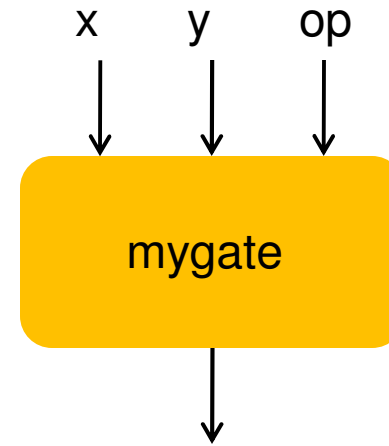
    function mygate;
        input x, y;
        input op;
        case (op)
            1'b0: mygate = x & y;
            1'b1: mygate = x | y;
            default: mygate = 1'b0;
        endcase
    endfunction

    assign q[0] = mygate(a, b, 0); // and gate
    assign q[1] = mygate(a, b, 1); // or gate

endmodule
```

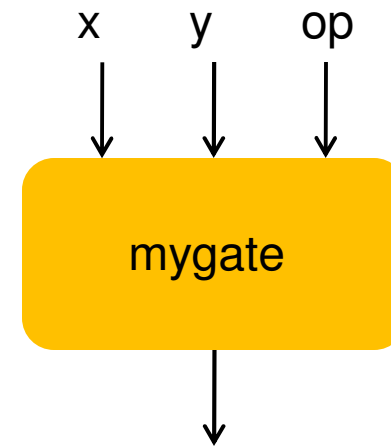
Functions

```
module andorgate(q, a, b);  
    output [1:0] q;  
    input      a, b;  
  
    function mygate;  
        input x, y;  
        input op;  
        case (op)  
            1'b0: mygate = x & y;  
            1'b1: mygate = x | y;  
            default: mygate = 1'b0;  
        endcase  
    endfunction  
  
    assign q[0] = mygate(a, b, 0); // and gate  
    assign q[1] = mygate(a, b, 1); // or gate  
  
endmodule
```



Functions

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module andorgate(q, a, b);  
    output [1:0] q;  
    input      a, b;  
  
    function mygate;  
        input x, y;  
        input op;  
        case (op)  
            1'b0: mygate = x & y;  
            1'b1: mygate = x | y;  
            default: mygate = 1'b0;  
        endcase  
    endfunction  
  
    assign q[0] = mygate(a, b, 0); // and gate  
    assign q[1] = mygate(a, b, 1); // or gate  
                x   y   op  
  
endmodule
```

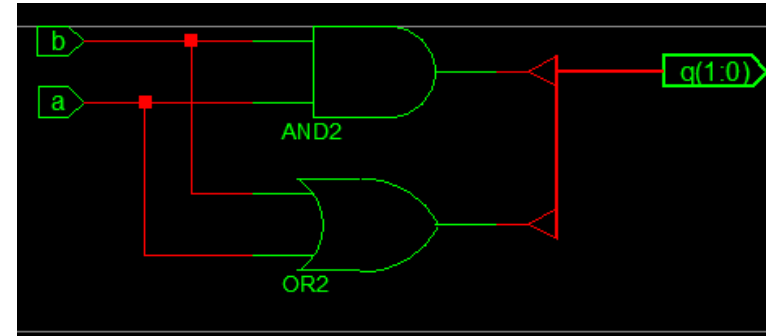


Functions

```
module andorgate(q, a, b);  
    output [1:0] q;  
    input      a, b;  
  
    function mygate;  
        input x, y;  
        input op;  
        case (op)  
            1'b0: mygate = x & y;  
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            default: mygate = 1'b0;  
        endcase  
    endfunction  
endmodule
```

```
assign q[0] = mygate(a, b, 0); // and gate  
assign q[1] = mygate(a, b, 1); // or gate
```

```
endmodule
```

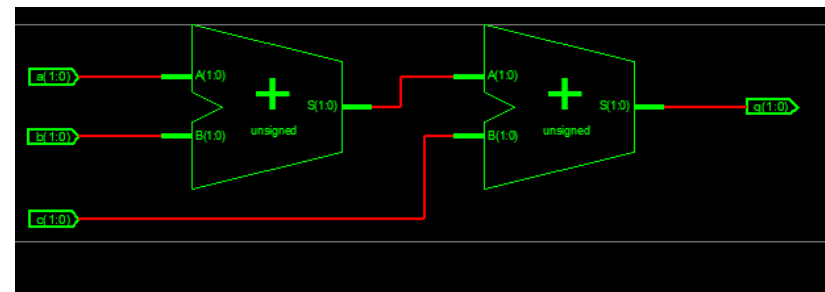


The logic is defined by the function calls, not by the function declaration

Functions can have local variables (reg)

```
module addit(q, a, b, c);  
    output [1:0] q;  
    input  [1:0] a, b, c;  
  
    function [1:0] myadd;  
        input [1:0] x, y, z;  
        reg [1:0] w;  
        begin  
            w      = x + y;  
            myadd = w + z;  
        end  
    endfunction  
  
    assign q = myadd(a, b, c);  
  
endmodule
```

← Two-bit output
← Local variable



Functions can have local variables (reg)

```
module addit(q, a, b, c);  
    output [1:0] q;  
    input  [1:0] a, b, c;  
  
    function [1:0] myadd;  
        input [1:0] x, y, z;  
        reg [1:0] w;  
        begin  
            w = x + y;  
            myadd = w + z;  
        end  
    endfunction  
  
    assign q = myadd(a, b, c);  
  
endmodule
```

← Two-bit output

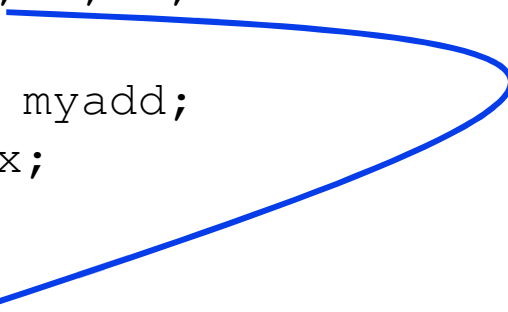
← Local variable

**w is shared across
all invocations of myadd**

**w is similar to a
static variable in C**

Functions can access module (global) variables

```
module addit(q, a, b, c);  
    output [1:0] q;  
    input  [1:0] a, b, c;  
  
    function [1:0] myadd;  
        input [1:0] x;  
        reg [1:0] w;  
        begin  
            w = a + x;  
            myadd = w + c;  
        end  
    endfunction  
  
    assign q = myadd(c); // will return c + a + c  
  
endmodule
```



Automatic Functions

```
module addit(q1, q2, a, b, c);  
  output [1:0] q1, q2;  
  input  [1:0] a, b, c;
```

```
  function automatic [1:0] myadd;  
    input [1:0] x, y, z;  
    reg [1:0] w;  
    begin  
      w      = x + y;  
      myadd = w + z;  
    end  
  endfunction
```

```
  assign q1 = myadd(a, b, c);  
  assign q2 = myadd(a, a, a);
```

```
endmodule
```

**automatic ensures that
all local variables
are truly local**

**With automatic, each
invocation of myadd will
use a different w**

Constant functions

- ❑ Don't create logic, but do compile-time calculations

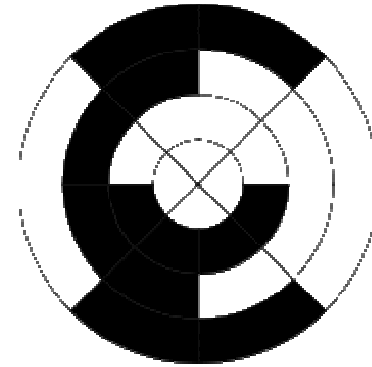
```
module ram(address);  
    parameter ram_depth = 256;  
    localparam w = bits(ram_depth);  
    input [w-1:0] address;  
  
    function bits;  
        input [31:0] v;  
        for (bits = 0; v > 1; bits = bits - 1)  
            v = v >> 1;  
    endfunction  
  
    // ..  
  
endmodule
```

Example: Gray to Bin Conversion Function

□ Remember Gray Code ?

- A gray code is a code where only 1 bit flips per code word

3-bit gray code	000
	001
	011
	010
	110
	111
	101
	100



□ How do we build a gray-to-bin function?

Example: Gray to Bin Conversion Function

- Gray to Bin should convert this:

IN	OUT
000	000
001	001
011	010
010	011
110	100
111	101
101	110
100	111

- It can be shown that

$$\text{bin-bit } i = \text{xor-sum}(\text{gray-bit } i \dots \text{gray-bit } N)$$

Example: Gray to Bin Conversion Function

- Gray to Bin should convert this:

IN	OUT	
000	000	
001	001	
011	010	
010	011	
110	100	bit3 = gray3 = 1
111	101	
101	110	
100	111	

- It can be shown that

$$\text{bin-bit } i = \text{xor-sum}(\text{gray-bit } i \dots \text{gray-bit } N)$$

Example: Gray to Bin Conversion Function

- Gray to Bin should convert this:

IN	OUT
000	000
001	001
011	010
010	011
110	100
111	101
101	110
100	111

$$\text{bit3} = \text{gray3} = 1$$

$$\text{bit2} = \text{gray3} \text{ xor } \text{gray2} = 1 \text{ xor } 1 = 0$$

- It can be shown that

$$\text{bin-bit } i = \text{xor-sum}(\text{gray-bit } i \dots \text{gray-bit } N)$$

Example: Gray to Bin Conversion Function

- Gray to Bin should convert this:

IN	OUT
000	000
001	001
011	010
010	011
110	100
111	101
101	110
100	111

$$\text{bit3} = \text{gray3} = 1$$

$$\text{bit2} = \text{gray3} \text{ xor } \text{gray2} = 1 \text{ xor } 1 = 0$$

$$\begin{aligned} \text{bit1} &= \text{gray3} \text{ xor } \text{gray2} \text{ xor } \text{gray1} \\ &= 1 \text{ xor } 1 \text{ xor } 0 = 0 \end{aligned}$$

- It can be shown that

$$\text{bin-bit } i = \text{xor-sum}(\text{gray-bit } i \dots \text{gray-bit } N)$$

Example: Gray to Bin Conversion Function

- Now let's build the gray to bin function

```
function [2:0] gray_to_bin;  
  input [2:0] a;  
  reg [2:0] q;  
  begin  
    q[2] = a[2];  
    q[1] = a[1] ^ a[2];  
    q[0] = a[0] ^ a[1] ^ a[2];  
    gray_to_bin = q;  
  end  
endfunction
```


Example: Gray to Bin Conversion Function

- ❑ Cool. Now let's use a loop to build the gray to bin function

```
function [W:0] gray_to_bin;  
    input [W:0] a;  
    reg [W:0] q;  
    integer i;  
    begin  
        for (i=0; i<=W; i = i + 1)  
            q[i] = ^(a >> i);  
        gray_to_bin = q;  
    end  
endfunction
```

Tasks

```
module operation;
    parameter delay = 5;
    reg [7:0] a, b, c;

    always @(a, b, c)
        runtask(a, b, c);    // same as: #5 a = b + c;

    task runtask;
    output [7:0] x;
    input [7:0] y, z;
    begin
        #delay x = y + z;
    end
endtask

    initial begin
        #5 b = 0; c = 2;
        #10 b = 8;
    end
endmodule
```

Tasks

```
module operation;
  parameter delay = 5;
  reg [7:0] a, b, c;

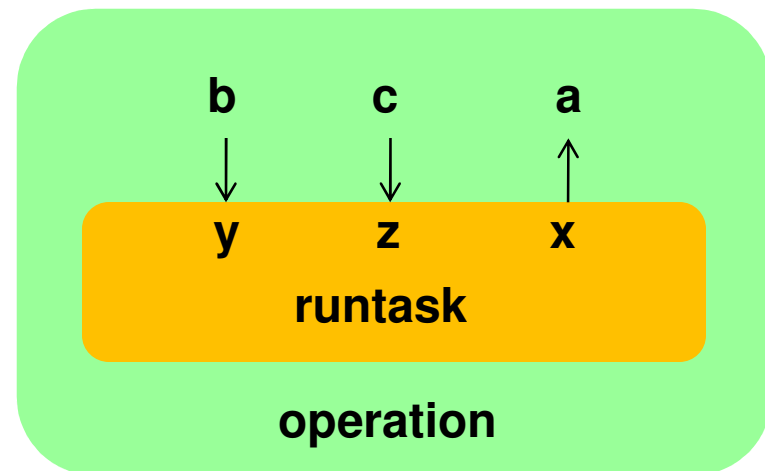
  always @(a, b, c)
    rntask(a, b, c); // same as: #5 a = b + c;

  task rntask;
  output [7:0] x;
  input [7:0] y, z;
  begin
    #delay x = y + z;
  end
endtask

initial begin
  #5 b = 0; c = 2;
  #10 b = 8;
end

endmodule
```

These are not inputs/outputs of the module, but links to the variables of the module



Tasks

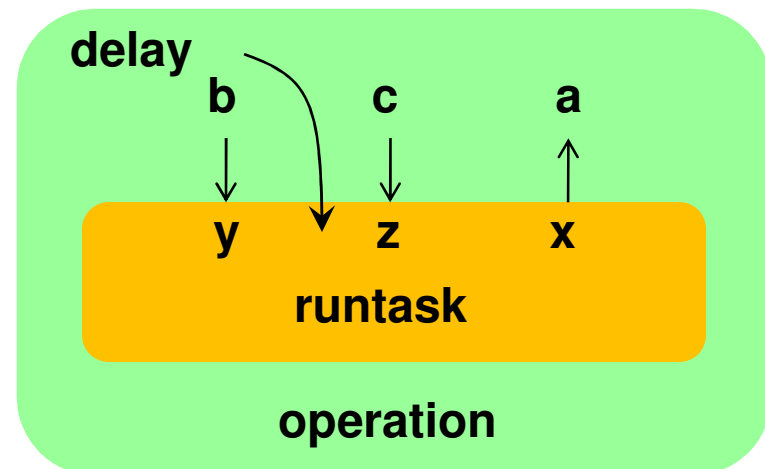
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  parameter delay = 5;
  reg [7:0] a, b, c;

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    runtask(a, b, c); // same as: #5 a = b + c;

  task runtask;
  output [7:0] x;
  input [7:0] y, z;
  begin
    #delay x = y + z;
  end
endtask

initial begin
  #5 b = 0; c = 2;
  #10 b = 8;
end
endmodule
```

Parameters, variables, etc
can be passed down from module
into a task



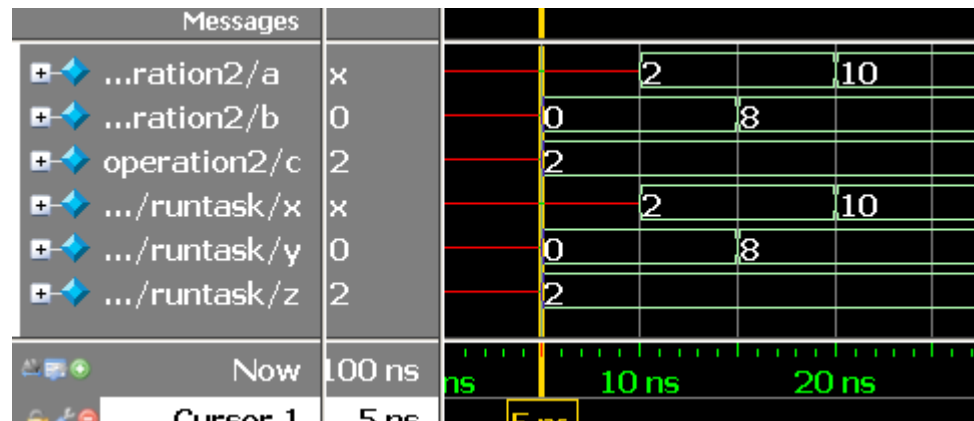
Tasks

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module operation;
  parameter delay = 5;
  reg [7:0] a, b, c;

  always @(a, b, c)
    runtask(a, b, c);    // same as: #5 a = b + c;

  task runtask;
  output [7:0] x;
  input [7:0] y, z;
  begin
    #delay x = y + z;
  end
  endtask

  initial begin
    #5 b = 0; c = 2;
    #10 b = 8;
  end
endmodule
```




Tasks

```
module operation;
  parameter delay = 5;
  reg [7:0] a, b, c;

  always @(a, b, c) begin
    runtask(a, b, c);
    #30 runtask(b, c, a);
  end
end
```

The moment when a task
is ready to run is called
enabling a task



```
task runtask;
  output [7:0] x;
  input [7:0] y, z;
  begin
    #delay x = y + z;
  end
endtask
```

```
initial begin
  #5 b = 0; c = 2;
  #10 b = 8;
end
```

```
endmodule
```

Tasks

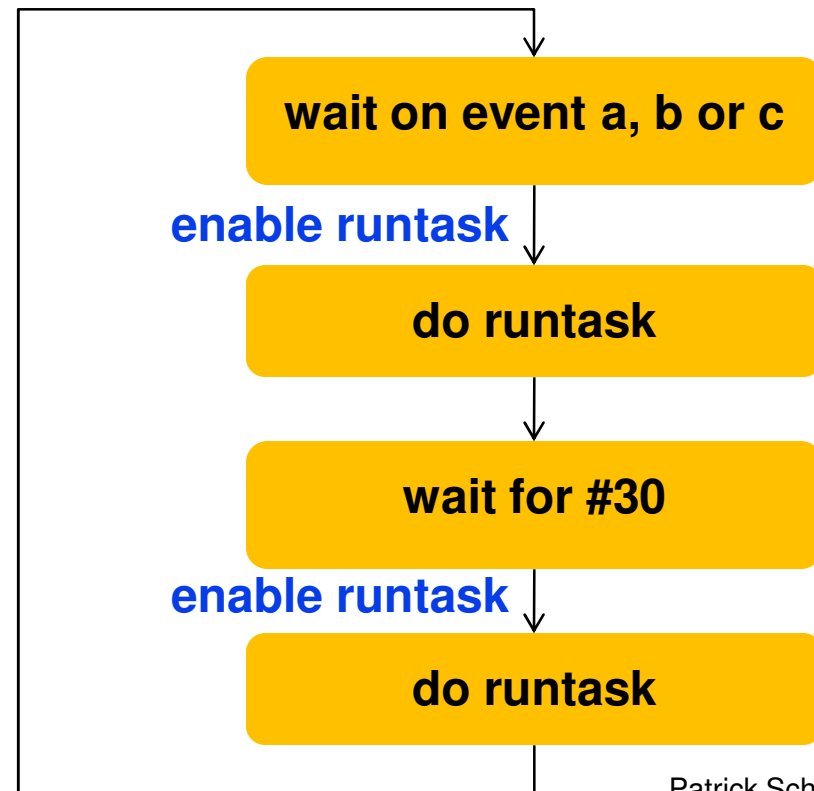
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module operation;  
  parameter delay = 5;  
  reg [7:0] a, b, c;  
  
  always @(a, b, c) begin  
    runtask(a, b, c);  
    #30 runtask(b, c, a);  
  end
```

```
task runtask;  
  output [7:0] x;  
  input [7:0] y, z;  
  begin  
    #delay x = y + z;  
  end  
endtask
```

```
initial begin  
  #5 b = 0; c = 2;  
  #10 b = 8;  
end
```

```
endmodule
```

The moment when a task is ready to run is called *enabling a task*



Functions and Tasks Summary

- ❑ Functions and Tasks enable shorthand notations within always and initial blocks
- ❑ Functions are single-output, and do not support event control
- ❑ Tasks may have more than a single output, and may have event-control statements

- ❑ In practice, only functions are used in synthesis
 - shorthand for combinational logic
 - calculation of compile-time constants